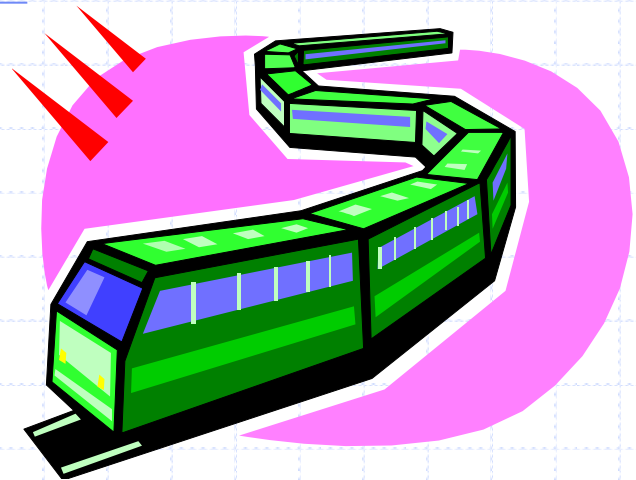


# Lists



# Position ADT

- ❑ The **Position** ADT models the notion of place within a data structure where a single object is stored
- ❑ It gives a unified view of diverse ways of storing data, such as
  - a cell of an array
  - a node of a linked list
- ❑ Just one method:
  - object **element()**: returns the element stored at the position

# Node List ADT

- The **Node List** ADT models a sequence of positions storing arbitrary objects
- It establishes a before/after relation between positions
- Generic methods:
  - **size()**, **isEmpty()**

## Accessor methods:

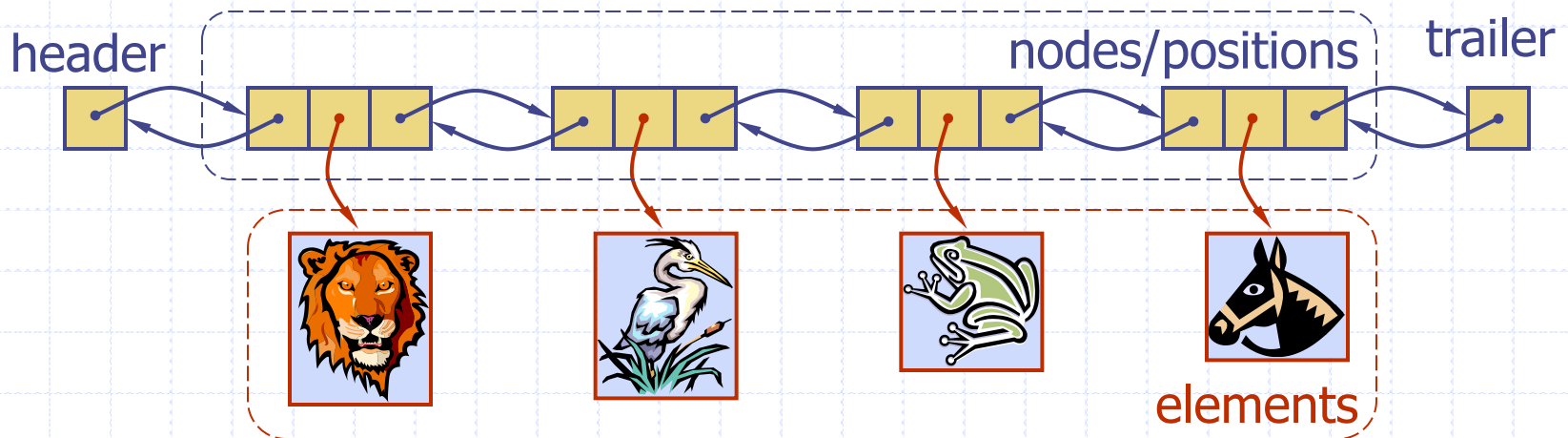
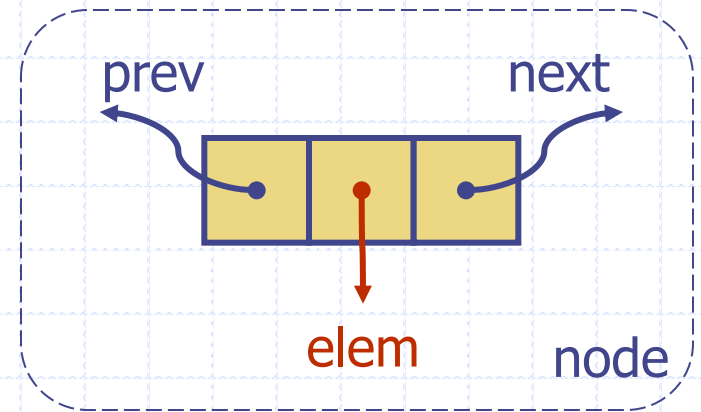
- **first()**, **last()**
- **prev(p)**, **next(p)**

## Update methods:

- **set(p, e)**
- **addBefore(p, e)**, **addAfter(p, e)**,
- **addFirst(e)**, **addLast(e)**
- **remove(p)**

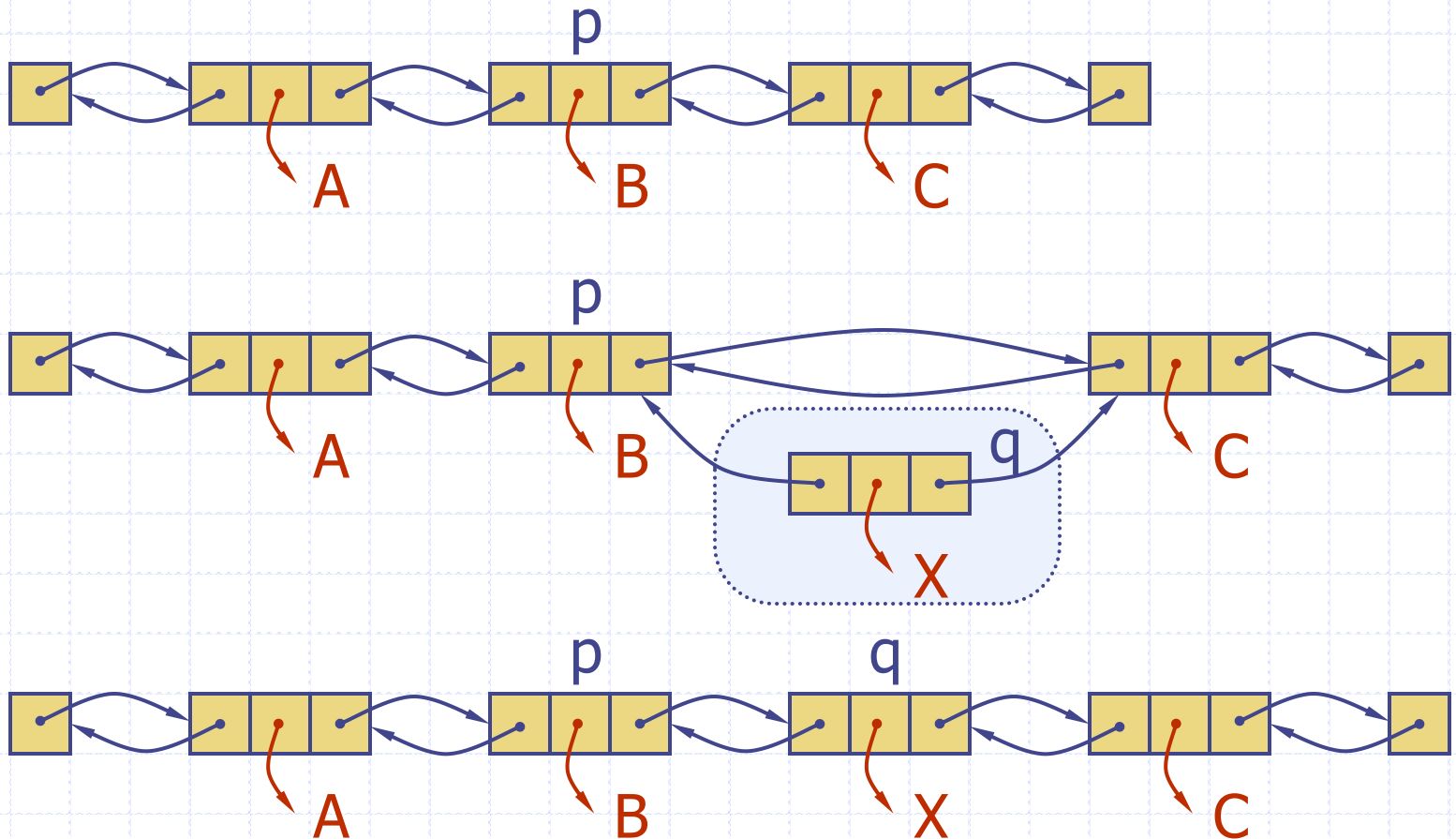
# Doubly Linked List

- A doubly linked list provides a natural implementation of the Node List ADT
- Nodes implement Position and store:
  - element
  - link to the previous node
  - link to the next node
- Special trailer and header nodes



# Insertion

- We visualize operation `insertAfter(p, X)`, which returns position `q`



# Insertion Algorithm

**Algorithm** `addAfter(p,e)`:

Create a new node `v`

`v.setElement(e)`

`v.setPrev(p)`      {link `v` to its predecessor}

`v.setNext(p.getNext())`    {link `v` to its successor}

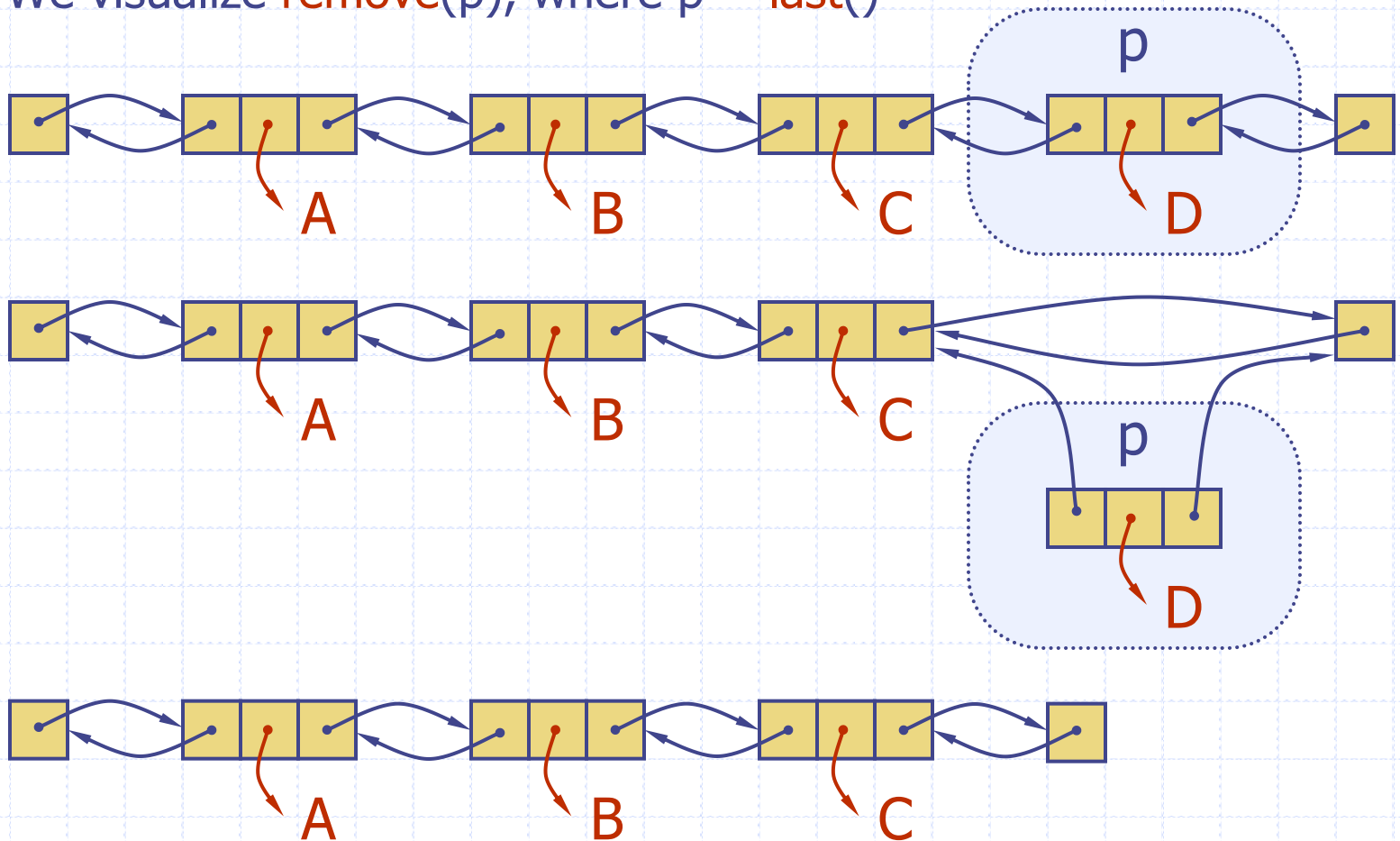
`(p.getNext()).setPrev(v)` {link `p`'s old successor to `v`}

`p.setNext(v)`      {link `p` to its new successor, `v`}

**return** `v`    {the position for the element `e`}

# Deletion

- We visualize `remove(p)`, where  $p = \text{last}()$



# Deletion Algorithm

**Algorithm** **remove**(p):

t = p.element      {a temporary variable to hold the  
return value}

(p.getPrev()).setNext(p.getNext())      {linking out p}

(p.getNext()).setPrev(p.getPrev())

p.setPrev(**null**)      {invalidating the position p}

p.setNext(**null**)

**return** t



# Performance

- In the implementation of the List ADT by means of a doubly linked list
  - The space used by a list with  $n$  elements is  $O(n)$
  - The space used by each position of the list is  $O(1)$
  - All the operations of the List ADT run in  $O(1)$  time
  - Operation `element()` of the Position ADT runs in  $O(1)$  time