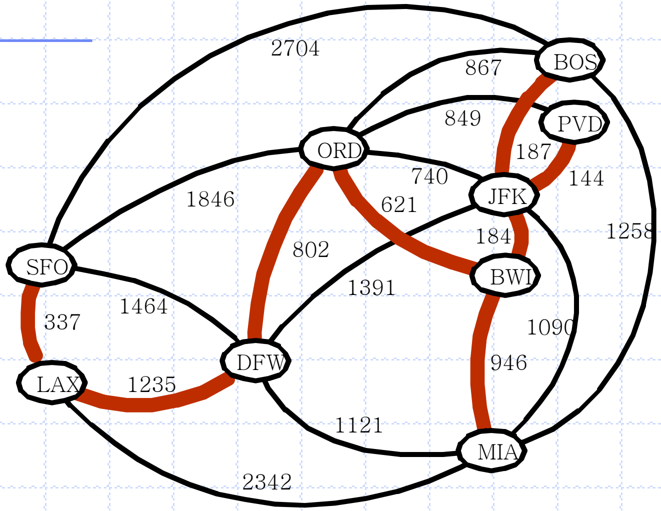
Minimum Spanning Trees



Minimum Spanning Trees

Spanning subgraph

- Subgraph of a graph G containing all the vertices of G

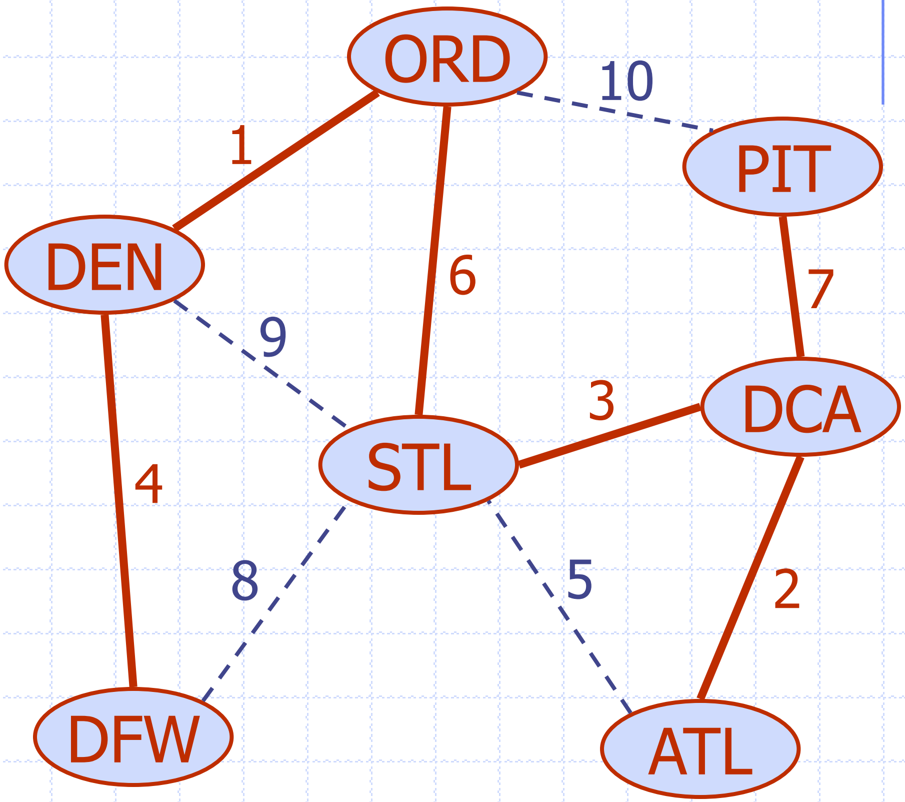
Spanning tree

- Spanning subgraph that is itself a (free) tree

Minimum spanning tree (MST)

- Spanning tree of a weighted graph with minimum total edge weight

- Applications
  - Communications networks
  - Transportation networks



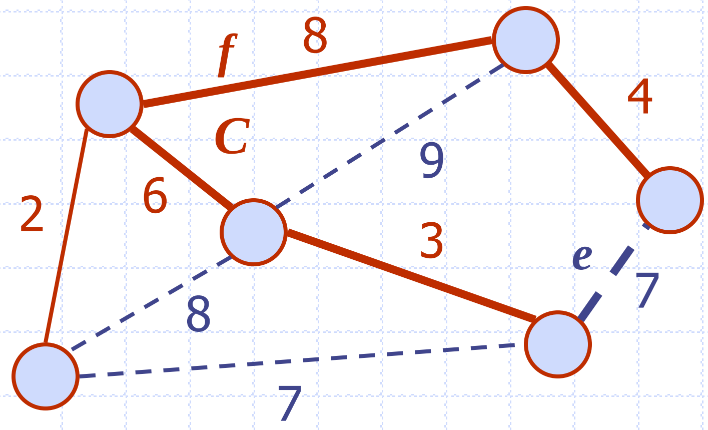
Cycle Property

Cycle Property:

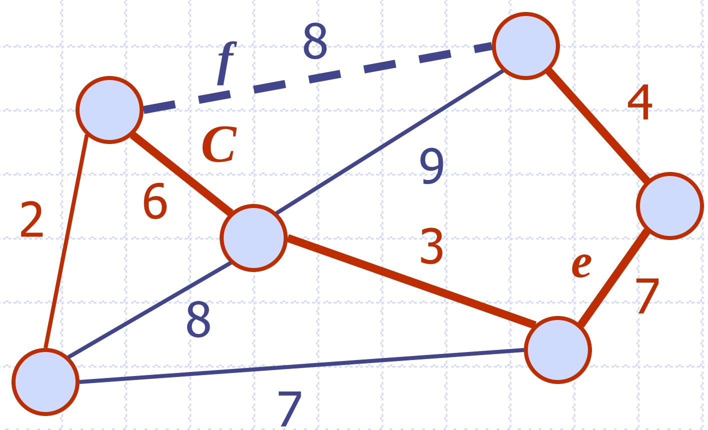
- Let T be a minimum spanning tree of a weighted graph G
- Let e be an edge of G that is not in T and C let be the cycle formed by e with T
- For every edge f of C, $weight(f) \leq weight(e)$

Proof:

- By contradiction
- If $weight(f) > weight(e)$ we can get a spanning tree of smaller weight by replacing e with f



Replacing f with e yields a better spanning tree



Partition Property

Partition Property:

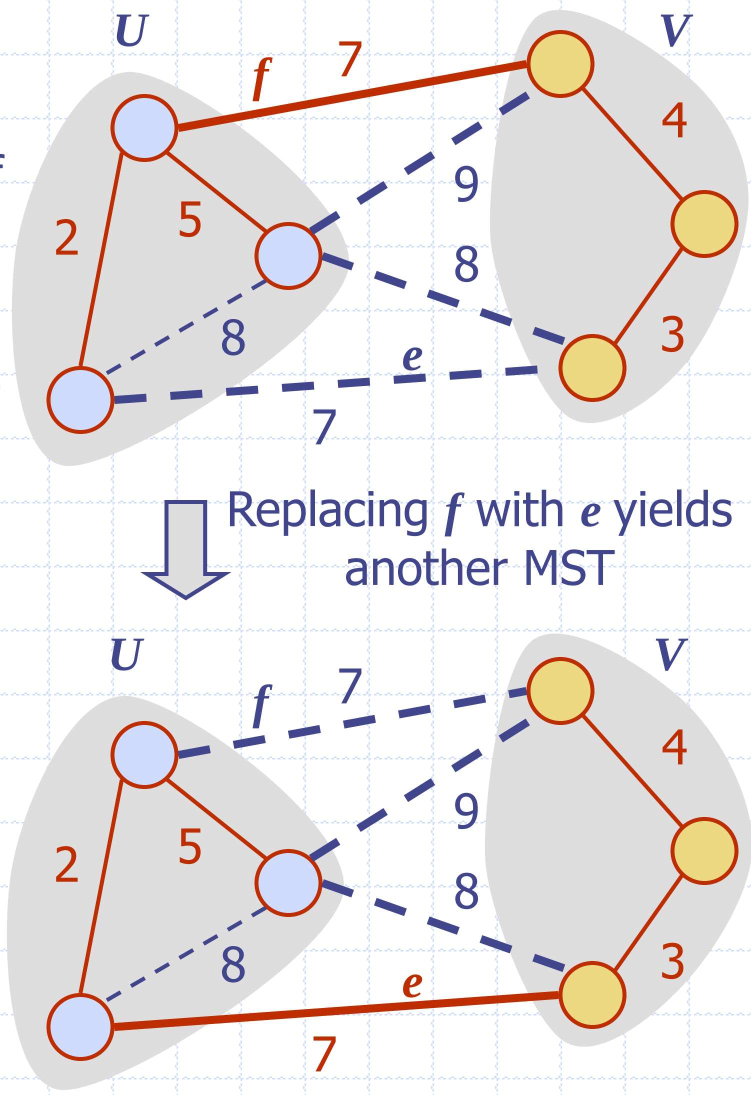
- Consider a partition of the vertices of G into subsets U and V
- Let e be an edge of minimum weight across the partition
- There is a minimum spanning tree of G containing edge e

Proof:

- Let T be an MST of G
- If T does not contain e, consider the cycle C formed by e with T and let f be an edge of C across the partition
- By the cycle property,

$$weight(f) \leq weight(e)$$

- Thus, $weight(f) = weight(e)$
- We obtain another MST by replacing f with e

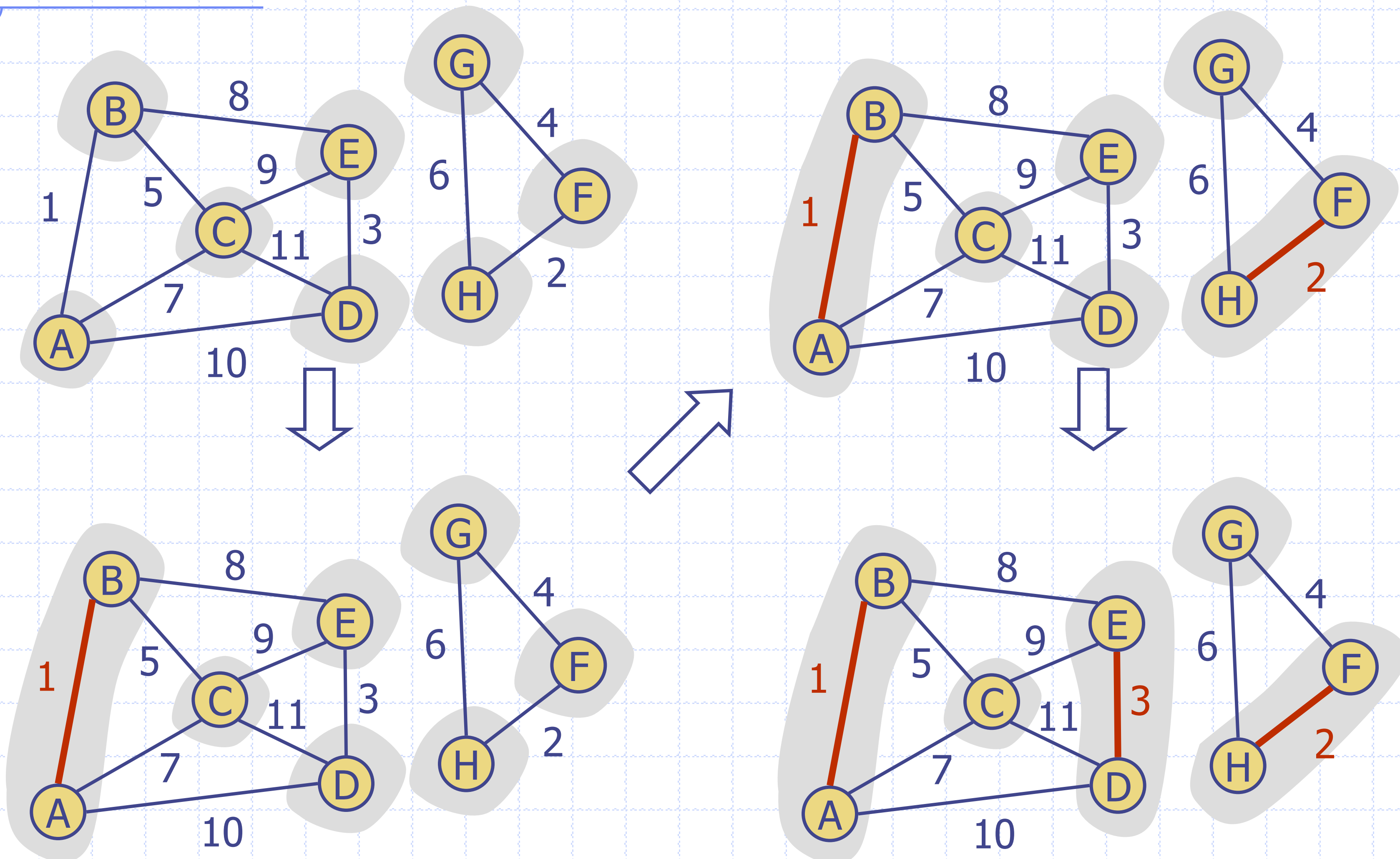


Kruskal's Algorithm

- Maintain a partition of the vertices into clusters
  - Initially, single-vertex clusters
  - Keep an MST for each cluster
  - Merge "closest" clusters and their MSTs
- A priority queue stores the edges outside clusters
  - Key: weight
  - Element: edge
- At the end of the algorithm
  - One cluster and one MST

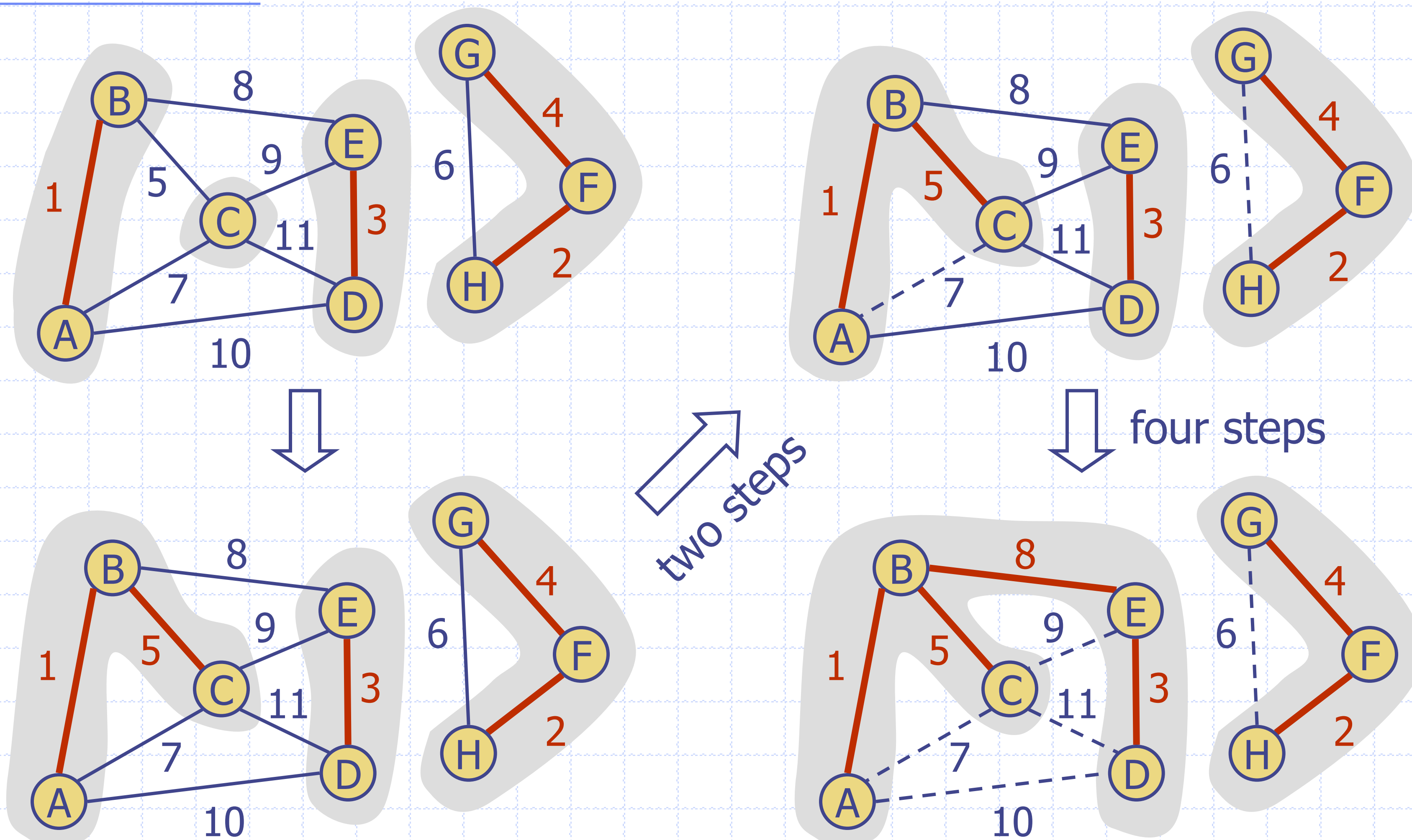
Algorithm ***KruskalMST*(*G*)**
  for each vertex ***v*** in ***G*** **do**
    Create a cluster consisting of ***v***
  let ***Q*** be a priority queue.
  Insert all edges into ***Q***
  $T \leftarrow \emptyset$
  {T is the union of the MSTs of the clusters}
  while T has fewer than $n - 1$ edges **do**
  $e \leftarrow Q.removeMin().getValue()$
    $[u, v] \leftarrow G.endVertices(e)$
    $A \leftarrow getCluster(u)$
    $B \leftarrow getCluster(v)$
    if $A \neq B$ **then**
      Add edge e to T
      $mergeClusters(A, B)$
  return T

Example



A B C D E F G H
1 5 7 8 9 10 11 3 6 4 2
A B C D E F G H
1 5 7 8 9 10 11 3 6 4 2
A B C D E F G H
1 5 7 8 9 10 11 3 6 4 2
A B C D E F G H
1 5 7 8 9 10 11 3 6 4 2

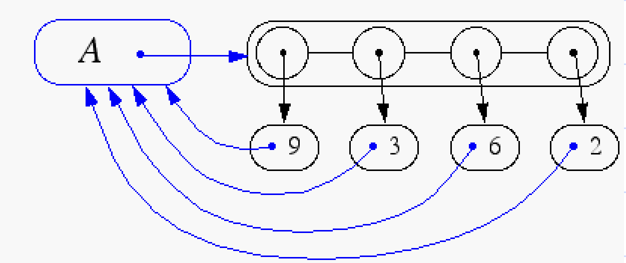
Example (contd.)



Data Structure for Kruskal's Algorithm

- The algorithm maintains a forest of trees
- A priority queue extracts the edges by increasing weight
- An edge is accepted it if connects distinct trees
- We need a data structure that maintains a partition, i.e., a collection of disjoint sets, with operations:
  - makeSet(u): create a set consisting of u
  - find(u): return the set storing u
  - union(A, B): replace sets A and B with their union

Recall of List-based Partition



- Each set is stored in a sequence
- Each element has a reference back to the set
  - operation find(u) takes O(1) time, and returns the set of which u is a member.
  - in operation union(A,B), we move the elements of the smaller set to the sequence of the larger set and update their references
  - the time for operation union(A,B) is min(|A|, |B|)
- Whenever an element is processed, it goes into a set of size at least double, hence each element is processed at most log n times

Partition-Based Implementation

- Partition-based version of Kruskal's Algorithm
  - Cluster merges as unions
  - Cluster locations as finds
- Running time $O((n + m) \log n)$
  - PQ operations $O(m \log n)$
  - UF operations $O(n \log n)$

Algorithm ***KruskalMST***(G)
  Initialize a partition P
  for each vertex v in G **do**
    $P.makeSet(v)$
  let Q be a priority queue.
  Insert all edges into Q
  $T \leftarrow \emptyset$
  {T is the union of the MSTs of the clusters}
  while T has fewer than $n - 1$ edges **do**
  $e \leftarrow Q.removeMin().getValue()$
    $[u, v] \leftarrow G.endVertices(e)$
    $A \leftarrow P.find(u)$
    $B \leftarrow P.find(v)$
    if $A \neq B$ **then**
      Add edge e to T
      $P.union(A, B)$
  return T

Prim-Jarnik's Algorithm

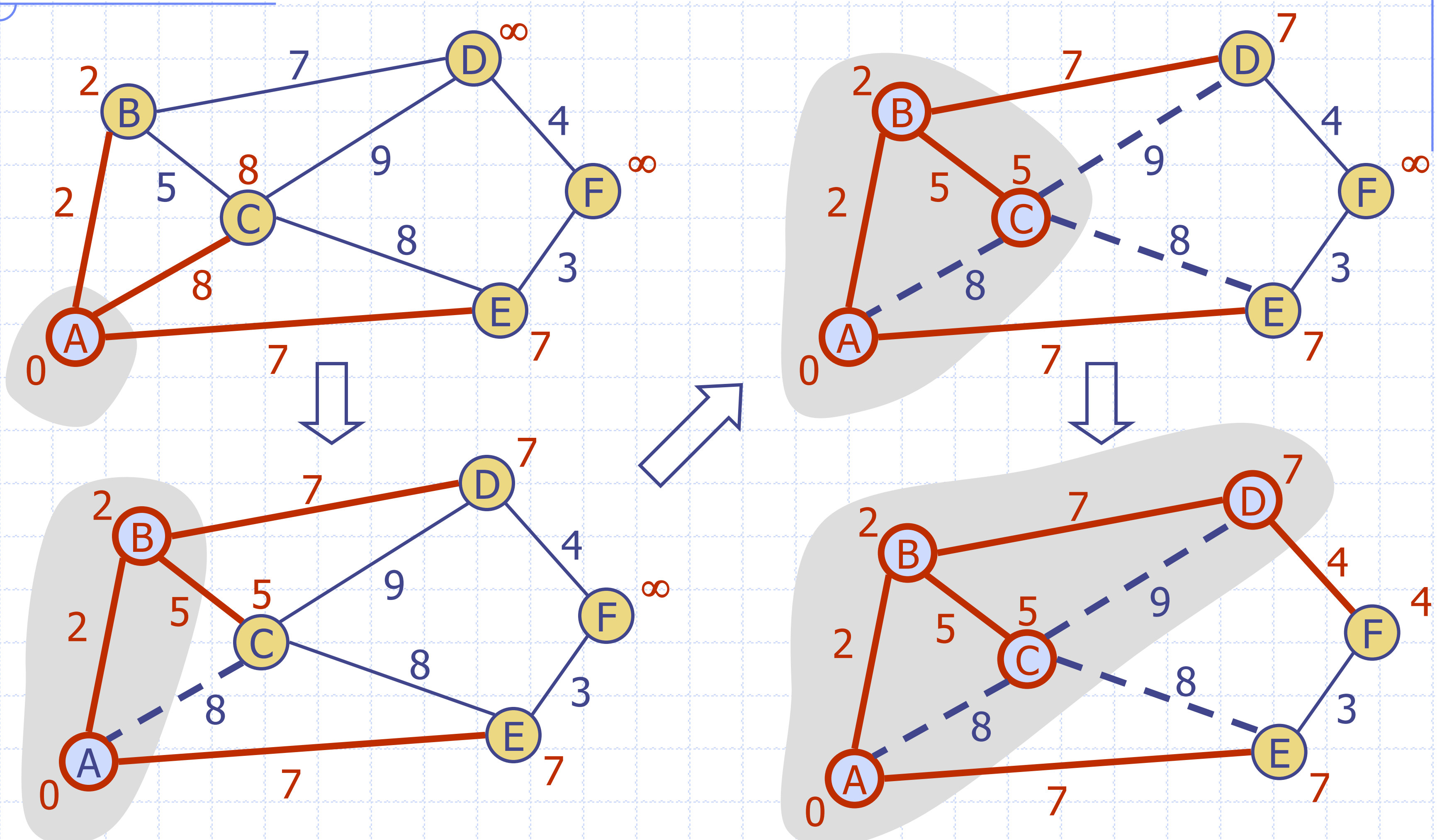
- Similar to Dijkstra's algorithm
- We pick an arbitrary vertex s and we grow the MST as a cloud of vertices, starting from s
- We store with each vertex v label $d(v)$ representing the smallest weight of an edge connecting v to a vertex in the cloud
- At each step:
  - We add to the cloud the vertex u outside the cloud with the smallest distance label
  - We update the labels of the vertices adjacent to u

Prim-Jarnik's Algorithm (cont.)

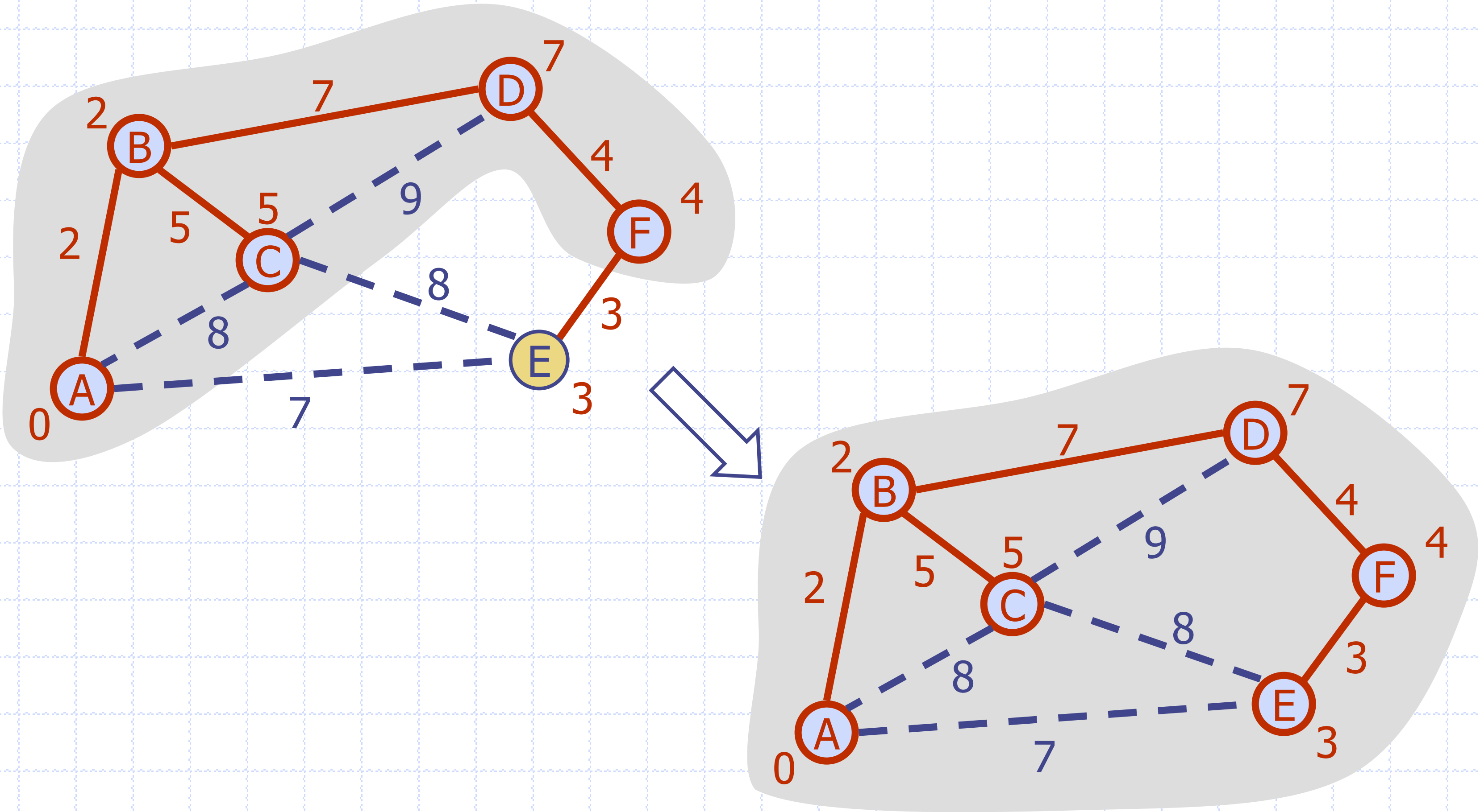
- A heap-based adaptable priority queue with location-aware entries stores the vertices outside the cloud
  - Key: distance
  - Value: vertex
  - Recall that method *replaceKey*(*l*,*k*) changes the key of entry *l*
- We store three labels with each vertex:
  - Distance
  - Parent edge in MST
  - Entry in priority queue

```
Algorithm PrimJarnikMST(G)
  Q ← new heap-based priority queue
  s ← a vertex of G
  for all v ∈ G.vertices()
    if v = s
      setDistance(v, 0)
    else
      setDistance(v, ∞)
    setParent(v, ∅)
    l ← Q.insert(getDistance(v), v)
    setLocator(v,l)
  while ¬Q.isEmpty()
    l ← Q.removeMin()
    u ← l.getValue()
    for all e ∈ G.incidentEdges(u)
      z ← G.opposite(u,e)
      r ← weight(e)
      if r < getDistance(z)
        setDistance(z, r)
        setParent(z,e)
        Q.replaceKey(getEntry(z), r)
```

Example



Example (contd.)



Analysis

- Graph operations
  - Method incidentEdges is called once for each vertex
- Label operations
  - We set/get the distance, parent and locator labels of vertex z $O(\deg(z))$ times
  - Setting/getting a label takes $O(1)$ time
- Priority queue operations
  - Each vertex is inserted once into and removed once from the priority queue, where each insertion or removal takes $O(\log n)$ time
  - The key of a vertex w in the priority queue is modified at most $\deg(w)$ times, where each key change takes $O(\log n)$ time
- Prim-Jarnik's algorithm runs in $O((n + m) \log n)$ time provided the graph is represented by the adjacency list structure
  - Recall that $\Sigma_v \deg(v) = 2m$
- The running time is $O(m \log n)$ since the graph is connected

Baruvka's Algorithm (Exercise)

- Like Kruskal's Algorithm, Baruvka's algorithm grows many clusters at once and maintains a forest T
- Each iteration of the while loop halves the number of connected components in forest T
- The running time is $O(m \log n)$

Algorithm ***BaruvkaMST***(G)
    $T \leftarrow V$ {just the vertices of G}
  while T has fewer than $n - 1$ edges **do**
    for each connected component C in T **do**
      Let edge e be the smallest-weight edge from C to another component in T
      if e is not already in T **then**
        Add edge e to T
  return T

Example of Baruvka's Algorithm (animated)

Slide by Matt Stallmann included with permission.